

Lecture 4, Jan 30, 2024

Addressing Modes

- *Inherent addressing*: address, data, or register is inherently specified by the instruction
 - Used by instructions that either don't take any operands, or operands are fixed by the instruction
 - e.g. NOP
- *Immediate addressing*: a constant value is explicitly stored within the code and used by the instruction
 - e.g. MOV R0, 0xAF
 - For word sizes above 1 byte, this will have to respect endianness
 - Best used to implement constants, since it is faster but not modifiable
- *Register addressing*: operands are registers
 - e.g. MOV R0, R1
 - Generally accumulator based or RISC architectures restrict where accumulators and general purpose registers can appear as operands
 - Best used to implement transfers or variable assignment
 - Limits us to the width of our registers
- *Direct addressing*: specifies a memory location to access (i.e. a constant pointer)
 - e.g. MOV [0xA000], 0xAF
 - The addresses themselves are stored within the instruction and are constant
 - Best used to implement global variables, constant loading, memory-mapped I/O, etc.
 - If we only use direct addressing, we need to make all variables global in scope and pre-assign fixed addresses to every variable
- *Indirect addressing*: the memory location in the operand contains the address of the data to be used by the instruction (i.e. a true pointer)
 - e.g. MOV R0, @[0xA000]
 - * The memory at 0xA000 is read, and interpreted as a memory address, and then data from that address is put into R0
 - Best used to implement pointers, arrays, etc, to implement data structures and pass data to subroutines
 - However, this uses 2 or more memory accesses per instruction, which are slow
- The below addressing modes are outside the minimum required set, so some microcontrollers might not support them, especially RISC
- *Register indirect addressing*: the register in the operand contains the address of the data to be used
 - e.g. MOV R0, @R1 or MOV R0, @@R1 (dereference twice, usually only on CISC)
 - Allows us to dynamically create and destroy new pointers
 - Allows doing math on pointers for e.g. array access with variable indices
 - Address space is again limited to register width, so we might not be able to access all the memory
 - Can become a vector for attack if a malicious actor can put data into the register
- *Indexed addressing*: a family of addressing modes to used to add indexing many of the previous modes
 - Relative addressing: relative to current program counter, e.g. BRR 0x05
 - * Useful for flow control
 - Base register: adding offset to register value, e.g. MOV R0, @[R1 + 5], MOV R0, @[R1 + R2]
 - * This is very useful for array dereferencing and data structures in general
- Other exotic modes (these are much less common)
 - Base plus index plus offset: effective address is the sum of all 3 (for supporting 2D arrays)
 - Scaled: effective address is the base address plus a scale times an index (for addressing arrays)
 - Register auto-increment/decrement

Example 1: Factorial

- Example: compute the factorial of a number stored in R0, assume 3 general purpose registers R0 to R2 are available and ignore checking pre-conditions
- Assume for the purposes of this example that we cannot modify R0

- Translate from C code:

```
int n;
int fact = 1;
for (int i = 1; i <= n; i++) {
    fact *= i;
}
```

- Create variable assignment table:

Variable	Assignment	Scope
n	R0	Global
fact	R1	Global
i	R2	Local to for

- Code generation:

```
MOV    R1, 0x01    ; int fact = 1; // note assumption that registers are 8-bit
MOV    R2, 0x01    ; int i = 1;
loop:
CMP    R2, R0
BGT    end        ; if (i > n) { break; }
MUL    R1, R2     ; fact = fact * i;
INC    R2         ; i++;
BRA    loop
end:
```

- Note we used labels here for the relative branch instructions; if not, we would have to compute the offsets ourselves
 - Assuming each instruction takes 2 bytes, the BGT would have an offset of 0x06 while the BRA would have 0xF6 (-10)

Example 2: Square Root

- Example: calculate the square root of a number S in R0 by the following procedure:

- Set $x_n = S/2$
- Iterate $x_{n+1} = \frac{1}{2} \left(x_n + \frac{S}{x_n} \right)$

- Code generation:

```
MOV    R1, R0     ; int x_n = S;
DIV    R1, 2      ; x_n = x_n / 2;
loop:
MOV    R2, R0     ; int temp = S;
DIV    R2, R1     ; temp = temp / x_n;
ADD    R1, R2     ; x_n = x_n + temp;
DIV    R1, 2      ; x_n = x_n / 2;
BRA    loop
```