## Lecture 5, Jan 17, 2024

## **Process Management**

- On Linux a process's state can be read through /proc/<PID>/status:
  - R: Running and runnable (running/waiting)
  - S: Interruptible sleep (blocked; can be resumed by the kernel if desired)
  - D: Uninterruptible sleep (blocked; cannot be resumed since it is waiting on I/O)
  - T: Stopped (can be continued explicitly by the user or another process)
  - Z: Zombie
- On Unix systems the kernel launches a single user process, init, which is the parent of all other processes
  - This is located at /sbin/init, and is usually systemd
  - This executes every other process on the machine and must always be active; if it exits the kernel will think you're shutting down
  - Some OSes will also create an idle process (e.g. Windows)

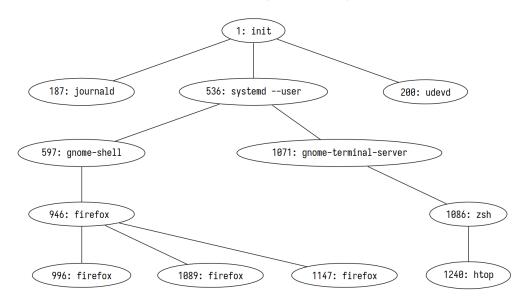


Figure 1: Example process tree.

- Each process is assigned a process ID (PID) on creation, which does not change and is unique for every active process
  - On most Linux systems this goes up to 32768; 0 is reserved/invalid
  - The OS can recycle a PID after the process dies
  - Each process has its own *address space* (i.e. its own copy of virtual memory)

## **Zombies and Orphans**

- The parent is responsible for the child and should acknowledge when the child terminates
- If the child exits first, it becomes a *zombie process* until the parent reads its exit status
  - The PCB cannot be removed by the OS until its exit status is read
  - Use the pid\_t wait(int \*status); syscall to check the child's status
    - \* Returns -1 on failure, 0 for nonblocking calls with no child changes, and the PID of the child on success
      - If there are multiple children, it returns the PID of the first child to terminate
    - \* The child's status is written to the address \*status, which is a bit mask
    - \* Use macros such as WIFEXITED(), WEXITSTATUS() etc to check specifics about the status

- wait() is a *blocking* system call, i.e. it will not return until the child is terminated
- The waitpid() syscall can be used to check on a child with a specific PID, and allows nonblocking calls
- When a child terminates the OS sends the parent a *signal* (SIGCHLD) to ask the parent to acknowledge the child
  - The parent is free to ignore this
  - If the parent always ignores it, the child will stay as a zombie until the parent dies, at which point it becomes an orphan and gets re-parented
- If the parent exits first, the child becomes an orphan process
  - Since some process still needs to acknowledge the child's exit, it needs a new parent
  - The OS will re-parent the child, usually to init
    - \* Note: A *subreaper* process (relatively new Linux feature) will take the place of **init** and adopt all orphans that are descendant from it