

Lecture 4, Jan 30, 2024

Constraint Satisfaction Problems (CSPs)

Definition

Constraint Satisfaction Problem: A CSP comprises of 3 components:

1. A set of variables $X = \{x_1, x_2, \dots, x_n\}$
2. A set of domains for each variable: $D: \{D_{x_1}, D_{x_2}, \dots, D_{x_n}\}$
3. A set of constraints C relating the variables

The problem is to find a value for each of the variables in its domain that satisfies all the constraints.

- Example: For the 4-queens problem:
 - Variables: $\{x_1, x_2, x_3, x_4\}$
 - Domain: for each variable: $D_{x_i} = \{1, 2, 3, 4\}$
 - Constraint: $\text{NoAttack}(x_i, x_j)$ (true if queen x_i can attack x_j)
 - * We can express this in a table, giving the value of NoAttack for every combination of x_i, x_j
- In general CSPs are NP-hard – no polynomial time solution exists
- But we can use heuristics to do better in a lot of the problems that arise in real life

Backtracking Search

- Assign each of the variables some value, and then check if the constraints are satisfied
- If the constraints are not satisfied, revert the last variable that we assigned – this is the process of *backtracking*
- If no value of the last variable works, then we go back one more variable and pick another value for that one, and so on
- This is essentially a brute force search if we pick values for the variables sequentially
 - However we can also use heuristics to aid our search

Algorithm 1 BacktrackingSearch(*prob, assign*)

```
1. if AllVarsAssigned(prob, assign) then
2.   if IsConsistent(assign) then
3.     return assign
4.   else
5.     return failure
6. var ← PickUnassignedVar(prob, assign)
7. for value ∈ OrderDomainValue(var, prob, assign) do
8.   assign ← assign ∪ (var = value)
9.   result ← BacktrackingSearch(prob, assign)
10.  if result ≠ failure then return result
11.  assign ← assign \ (var=value)
12. return failure
```

Figure 1: Basic backtracking search algorithm.

- The pseudocode above is a template for the backtracking search algorithm
 - Each level of the recursion picks a variable to set, goes through all values of that variable and checks if any of them work
 - We can specify different implementations for $\text{PickUnassignedVariable}$ and OrderDomain
- One simple improvement we can make is to only assign variables to values that satisfy all the constraints

- Otherwise we would do a lot of meaningless searches when we pick a value that violates a constraint and keeps assigning others
- e.g. for N -queens, check that the new queen cannot attack any previous queens before placing it

Algorithm 2 BacktrackingSearch($prob, assign$)

```

1. if AllVarsAssigned( $prob, assign$ ) then
2.  $var \leftarrow$  PickUnassignedVar( $prob, assign$ )
3. for  $value \in$  OrderDomainValue( $var, prob, assign$ ) do
4.   if VallsConsistentWithAssignment( $value, assign$ ) then
5.      $assign \leftarrow assign \cup (var = value)$ 
6.      $result \leftarrow$  BacktrackingSearch( $prob, assign$ )
7.     if  $result \neq failure$  then return result
8.      $assign \leftarrow assign \setminus (var=value)$ 
9. return failure

```

Figure 2: Improved backtracking search algorithm.

- Can we do better and reduce our backtracks even more?
 - Every time we assign a variable, it reduces the domain that the other variables can take based on constraints
 - The domain reduction (restrictions) on the other variables can propagate to even more variables
 - Every time we do an assignment, we call the inference function, which restricts the domain further based on the constraints and the new variable value
 - * This can look at multiple constraints at the same time but often we stick to just 1

Algorithm 3 BacktrackingSearch_with_Inference($prob, assign$)

```

1. if AllVarsAssigned( $prob, assign$ ) then
2.  $var \leftarrow$  PickUnassignedVar( $prob, assign$ )
3. for  $value \in$  OrderDomainValue( $var, prob, assign$ ) do
4.   if VallsConsistentWithAssignment( $value, assign$ ) then
5.      $assign \leftarrow assign \cup (var = value)$ 
6.      $inference \leftarrow$  Infer( $var, prob, assign$ )
7.     if  $inference \neq failure$  then
8.        $assign \leftarrow assign \cup inference$ 
9.        $result \leftarrow$  BacktrackingSearch( $prob, assign$ )
10.    if  $result \neq failure$  then return result
11.     $assign \leftarrow assign \setminus ((var=value) \cup inference)$ 
12. return Failure

```

Figure 3: Improved backtracking with inference.